This is a closed-book exam. You may use two double-sided 8.5x11 sheets of notes.

All answers to this exam should be written in your exam booklet(s). Start with the questions that you know how to do, and try not to spend too long on any one question. Partial credit will be granted where appropriate. You will have two hours and twenty minutes. Good luck!

1. Vocabulary. (16 points) Identify the term defined in class that matches the following definitions. If the answer is an acronym, specify what it stands for.

   a. A cache that stores the page table entry for recently accessed pages.

   b. A page replacement algorithm that measures how much pages are used and replaces ones that aren’t used much.

   c. A rule of thumb suggesting that for a memory system of variable-sized partitions in equilibrium, there should be twice as many partitions as holes.

   d. This piece of hardware attached to the CPU translates virtual memory names into physical addresses.

   e. This is the name for the process whereby an I/O device transfers data to or from the main memory, bypassing the CPU.

   f. This principle or property of an executing program suggests that looking at the behavior in the recent past will indicate the likely behavior in the near future.

   g. This term refers to a set of disk sectors that are simultaneously accessible without moving the read/write head, because they are located on corresponding tracks on each platter.

   h. A software construct, defined to have nonnegative integral values, with two possible operations, sometimes called wait and signal.

2. Process State. (8 points) A particular operating system will have five process states: RUNNING, READY, BLOCKED, BLOCKED/SUSPEND, and READY/SUSPEND. (The latter two states apply to processes whose working set is not currently in memory.) Draw a transition diagram showing the allowable transitions between states. Label each transition with the mechanism that would cause that transition to take place.
3. **Race Conditions.** (12 points) For each of the following programs, state the possible values for \( x \) and \( y \) if the program terminates. Express your answer as a set of \((x,y)\) pairs. For example, the starting configuration should be written as \{(8,3)\}.

In addition, state whether the program will always terminate, sometimes terminate, or never terminate.

You may assume that the variables are initialized as follows:

- \( x \): integer **init** 8
- \( y \): integer **init** 3
- \( m0 \): general semaphore **init** 0
- \( m1 \): general semaphore **init** 1

a). \texttt{cobegin} \( x := x-y; \ // \ y := y-x; \ \texttt{coend} \)

b). \texttt{cobegin}
\[
\text{DOWN}(m1); \ x := x-y; \ \text{UP}(m1);
\ //
\text{DOWN}(m1); \ y := y-x; \ \text{UP}(m1);
\ \texttt{coend}
\]

c). \texttt{cobegin}
\[
\text{DOWN}(m0); \ x := x-y; \ \text{UP}(m0);
\ //
\text{DOWN}(m0); \ y := y-x; \ \text{UP}(m0);
\ \texttt{coend}
\]

d). \texttt{cobegin}
\[
\text{DOWN}(m1); \ x := x-y; \ \text{UP}(m1); \ \text{UP}(m0);
\ //
\text{DOWN}(m0); \ y := y-x; \ \text{UP}(m0); \ \text{UP}(m1);
\ \texttt{coend}
\]

4. **Paging and Latency.** (8 points) A large employee database is implemented on a system with 64-bit virtual memory. As a result, the employee records can be stored within the virtual address space as an array of records, in order according to the employee ID number. A record may be accessed by its ID number, or by social security number (SSN). Both access methods use hash tables and are similar in speed. Your boss wants you to write a function that will compute the total salary for all workers. Does it matter which method you use for access? Why or why not? (Assume that employee ID numbers are based upon the hire date and are thus unrelated to SSN, and that several records fit onto a page in memory.)

5. **Linux Processes** (8 points). Explain how to set and/or modify the priority of a process in Linux, assuming the kernel is using the `SCHED_OTHER` protocol. What limitations are placed on ordinary users (as opposed to root) as far as the ability to set process priorities?
6. **Virtual Memory.** (18 points) You’re working on the memory manager for a new computer, code-named Elephant. Elephant will use a pure segmented memory system, with up to eight segments available. The segment to be used will be indicated by the first three bits of the virtual address. Segments are allocated in sizes that are multiples of 1K, and every segment must begin in memory at an address that is a multiple of 1K. The base and limit of the segment may then be given in the segment table with the last ten bits left off.

Below is a sample segmentation table for Elephant. For the memory actions given below, specify (i) which faults, if any, would occur as a result; (ii) specify the address in memory that would come out of the MMU, if any; (iii) specify any changes to the segment table that would occur, if any.

<table>
<thead>
<tr>
<th>Segment</th>
<th>Permissions</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>Resident</td>
</tr>
<tr>
<td>000:</td>
<td>1</td>
</tr>
<tr>
<td>001:</td>
<td>1</td>
</tr>
<tr>
<td>010:</td>
<td>1</td>
</tr>
<tr>
<td>011:</td>
<td>1</td>
</tr>
<tr>
<td>100:</td>
<td>1</td>
</tr>
<tr>
<td>101:</td>
<td>1</td>
</tr>
<tr>
<td>110:</td>
<td>0</td>
</tr>
<tr>
<td>111:</td>
<td>0</td>
</tr>
</tbody>
</table>

a. Read operation to address 00110101010101010.
b. Read operation to address 11010000000001010.
c. Read operation to address 00000000000000010.
d. Write operation to address 00110101010101010.
e. Write operation to address 10011010101010101.
f. Write operation to address 01100101010101010.

7. **Page Replacement.** (8 points) Slick Willy has proposed to install a new page replacement algorithm on your computer. He shows you the results of an experiment he did, where he counts the page faults generated by a particular process allocated a certain number of frames to run. The count shows the algorithm does nearly as well as OPT, and much better than LRU on this experiment. Assuming you believe the results of this test are true, what concerns might keep you from adopting the new policy despite such impressive results? (You should be able to give at least two separate considerations.)
8. **Bit Calculations.** (6 points) A newly designed operating system will use a segmented, paged virtual memory system (similar to the Multics virtual memory) with 32-bit virtual addresses. Each process will have access to up to 64 segments, and each segment will have its own page table.

   a. How many bits are available for the page name and page offset combined? What is the maximum length of a segment, in words?

   b. Suppose that pages are $2^f$ words in size. In terms of $f$, what is the expected (average) internal fragmentation per segment?

   c. What is the expected table fragmentation per segment, again in terms of $f$? Assume that each page takes one word to describe in the page table.

9. **Ext2 File System** (8 points). The diagram below shows the data structures maintained within a block group of the Ext2 file system. Identify all the data structures that would need to be used in order to complete each of the following tasks. Assume that nothing is cached in memory, so that all information for each part must be read in from the disk. (Note: each part will typically use several of the data structures below.)

```
        SUPER BLOCK | GROUP DESCRIPTORS | BLOCK DESCRIBER | INODE TABLE | DATA BLOCKS
```

a). Determine how much free space is available on the disk.

b). Locate a free inode and mark it used.

c). Check the modification date of the root directory.

d). Read the contents of the root directory of the disk.

10. **Semaphores and Deadlock.** (8 points) Consider the following process definitions involving semaphores. Which set(s) of processes might become deadlocked? Draw resource allocation diagrams for any scenarios where deadlock is possible.

```
init:   semaphore s1(1), s2(1), s3(1);

proc1: process
   DOWN(s1); DOWN(s2); DOWN(s3); UP(s3); UP(s2); UP(s1);
endprocess;

proc2: process
   DOWN(s1); DOWN(s3); UP(s3); UP(s1);
endprocess;

proc3: process
   DOWN(s2); DOWN(s1); UP(s1); UP(s2);
endprocess;
```